(Re)writing your SQL for optimal performance



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Some projects I have worked on

- Italian Chambers of Commerce:
 - Reporting engine (Visure/Certificati)
 - Replication solution Oracle/Unix to DB2/MVS
- 2000: one of the first browser based business application development tools
- Java ERP development framework
- NebulaERP, italian ERP software
- Database Events notification system
- heterogeneous, over the Internet, data replication solution
- IoT: PostgreSQL based Cloud API server

SQL Performance facets

- Database design
 - Logical, physical (indexing)
- Application design
 - How you interact with the rdbms (ORM?!)
 - Functionality used (e.g. prepared statements, batching...)
- Database tuning and management (e.g. statistics, vacuum), O/S tuning
- Optimizer related issues

SQL is a declarative language

- RDBMS reads your code, parses it and feeds the result to the optimizer
- The optimizer figures out the best plan to execute your query
- SQL optimization is a NP-complete problem
- Optimizers are imperfect so sometimes they return (slow) sub-optimal plans

Manual SQL optimization

- Some RDBMS allow hints to be fed to the planner
- PostgreSQL does not implement hints (proudly)
- BUT, most of the times, we can work around the "issue" by rewriting the SQL
- We will see two real world examples

Reading a PostgreSQL plan

- A plan is a tree shaped structure
- Leaves correspond to data access operations (tables and indexes)
- Intermediate nodes correspond to operators applied to child nodes (e.g. join, sort...)
- To visualize plans I will use PEV from Alex Tatiyants, a handy tool for the job
- Use EXPLAIN (ANALYZE, COSTS, VERBOSE, BUFFERS, FORMAT JSON) to generate JSON plan description
- Paste result into browser based UI

Example 1: The original SQL

```
CREATE or replace VIEW saldi sottoconti AS
                                                                        , rc.codice part
SELECT rc.codice azienda
    , CASE
                                                                       , rc.codice template fiscale
       WHEN rc.id divisione IS NULL
                                                                       , rc.id pdc
      THEN '-'
                                                                       , rc.id sottoconto
      ELSE rc.id divisione::text
     END || '#' || d.id esercizio comp::text || '#' || rc.id sottoconto::
                                                                        , s.tree path
    , rc.id divisione
                                                                          , cee.tree path || '|' || s.codice as tree path cee
    , d.id esercizio comp
                                                                  FROM documenti d
    , sum (CASE
          WHEN rc.seqno = 1
                                                                  join tipi doc az tda
           AND rc.extra contabile = '0'
                                                                   on tda.codice azienda = d.codice azienda
          AND coalesce(tda.sottotipo, '-') <> 'CH'
                                                                   and tda.codice tipo doc = d.codice tipo doc
      THEN rc.importo
          ELSE 0
                                                                   and tda.codice tipo doc az = d.codice tipo doc az
         END) AS totale dare
                                                                  JOIN reg contabili rc
    , sum(CASE
                                                                    ON d.codice azienda = rc.codice azienda
          WHEN rc.seqno = (-1)
           AND rc.extra contabile = '0'
                                                                   AND d.id documento = rc.id documento
           AND coalesce(tda.sottotipo, '-') <> 'CH'
                                                                  JOIN sottoconti s
      THEN rc.importo
                                                                    ON rc.codice part = s.codice part
          ELSE 0
         END) AS totale avere
                                                                   AND rc.codice template fiscale = s.codice template fiscale
    , sum (CASE
                                                                   AND rc.id pdc = s.id pdc
          WHEN rc.segno = 1
                                                                   AND rc.id sottoconto = s.id sottoconto
           AND rc.extra contabile = '0'
          AND coalesce(tda.sottotipo, '-') = 'CH'
                                                                  JOIN conti c
      THEN rc.importo
                                                                    ON s.codice part = c.codice part
          ELSE 0
                                                                   AND s.codice template fiscale = c.codice template fiscale
         END) AS totale dare ch
    , sum (CASE
                                                                   AND s.id pdc = c.id pdc
          WHEN rc.segno = (-1)
                                                                   AND s.id conto = c.id conto
           AND rc.extra contabile = '0'
                                                                  --left join cee cee
           AND coalesce(tda.sottotipo, '-') = 'CH'
      THEN rc.importo
                                                                            on cee.id cee = s.id cee
           ELSE 0
                                                                  WHERE d.stato = 'C'
         END) AS totale avere ch
                                                                  GROUP BY rc.codice azienda
    , sum (CASE
          WHEN rc.seqno = 1
                                                                          , d.id esercizio comp
           AND rc.extra contabile = '1'
                                                                          , c.tipo mastro
          AND coalesce(tda.sottotipo, '-') <> 'CH'
                                                                          , rc.id divisione
      THEN rc.importo
          ELSE 0
                                                                          , s.codice
         END) AS totale dare ec
                                                                          , rc.codice part
    , sum (CASE
                                                                          , rc.codice template fiscale
          WHEN rc.seqno = (-1)
           AND rc.extra contabile = '1'
                                                                          , rc.id pdc
           AND coalesce(tda.sottotipo, '-') <> 'CH'
                                                                          , rc.id sottoconto
       THEN rc.importo
                                                                          , s.descrizione
          ELSE 0
         END) AS totale avere ec
                                                                          , s.tree path
```

The query

```
SELECT *
FROM saldi_sottoconti
WHERE
     codice_azienda = 'S0040SC'
AND id_esercizio_comp = 2017
AND id_divisione = 2
```

The (slow) plan



Have a look at the slowest node

	INDEX SCAN on solving.sottoconti (s) using soco_pk	8.59s 79 %
93ms 5 %	Index Scan Node finds I perform 2 read operations actual value from the table	relevant records based on an Index . Index Scans :: one to read the index and another to read the e.
	Node Type	Index Scan
	Parent Relationship	Inner
	Parallel Aware	false
	Scan Direction	Forward
	Index Name	soco_pk
	Relation Name	sottoconti
	Schema	solving
	Alias	s
	Startup Cost	0.42
	Total Cost	1.59
	Plan Rows	1
	Plan Width	88
	Actual Startup Time	0.007
	Actual Total Time	0.007
	Actual Rows	1
	Actual Loops	1226906

s.id_sottoconto,s.codice,s.descrizione,s.i

Index scan should be fast and it is (0.007 ms) But this node is executed 1,226,906 times (See Actual Loops)

So the problem is caused by the left node in the join operation that is feeding too much data into this node

This node generates 1.2 million rows

largest bad estimate	
Nested Loop Node mer ecord in the first set and in natching records are retu	ges two record sets by looping through every trying to find a match in the second set. All Irned.
Node Type	Nested Loop
Parent Relationship	Outer
Parallel Aware	false
Join Type	Inner
Startup Cost	5.95
Total Cost	4054.04
Plan Rows	80
Plan Width	88
Actual Startup Time	0.105
Actual Total Time	1019.141
Actual Rows	1226906
Actual Loops	1
Output	d.id_esercizio_comp,tda.sottotipo,rc.codi ce_azienda,rc.id_divisione,rc.id_sottocont o,rc.segno,rc.extra_contabile,rc.importo, rc.codice_part,rc.codice_template_fiscal e,rc.id_pdc,c.tipo_mastro,c.codice_part,c. codice_template_fiscale,c.id_pdc,c.id_con to
Shared Hit Blocks	144466
Channel David Diastra	0

So we descend to the children nodes to further investigate

Here are the children

	NESTED LOOP	964.5ms 9 %			INDEX SCAN	8.59s 7 9
	Inner join largest bod estimate				on solving.sottoconti (s) using soco_pk	
			THEFY COAN			
aper join	41.41ms 0 %		INDEX SCAN	546.93ms 5 %		
costilest	=		using cnti_pk	=		
Nested Loop Node merg ecord in the first set and t matching records are retu	ges two record sets by looping through every rying to find a match in the second set. All rned.		Index Scan Node finds rele perform 2 read operations: or actual value from the table.	vant records based on an Index . Index Scans ne to read the index and another to read the		
Node Type	Nested Loop		Node Type	Index Scan		
Parent Relationship	Outer		Parent Relationship	Inner		
Parallel Aware	false		Parallel Aware	false		
Join Type	Inner		Scan Direction	Forward		
Startup Cost	5.67		Index Name	cnti_pk		
Total Cost	3986.72		Relation Name	conti		
Plan Rows	187		Schema	solving		
Plan Width	59		Alias	c		
Actual Startup Time	0.079		Startup Cost	0.29		
Actual Total Time	54.604		Total Cost	0.35		
Actual Rows	14782		Plan Rows	1		
Actual Loops	1		Plan Width	29		
	d.id_esercizio_comp,tda.sottotipo,rc.codi		Actual Startup Time	0.019		
Output	ce_azienda,rc.id_divisione,rc.id_sottocont o,rc.segno,rc.extra_contabile,rc.importo,		Actual Total Time	0.037		
	rc.codice_part,rc.codice_template_fiscal e.rc.id_pdc		Actual Rows	83		
Shared Hit Blocks	26210		Actual Loops	14782		
Shared Read Blocks	0			c.id_conto,c.codice,c.descrizione,c.id_con to_parent,c.stato,c.tipo_mastro,c.ca,c.cli		
Shared Dirtied Blocks	0			for,c.id_cee,c.id_riclass,c.tree_path,c.codi ce_template_fiscale,c.id_pdc,c.codice_azi		
Shared Written Blocks	0		Output	enda,c.codice_tipo_sottoconto_anag,c.fla		
Local Hit Blocks	0			empre_sottoconti,c.tipo_cespite,c.codice stampa c flag bide codice c decesizione		
Local Read Blocks	Local Read Blocks 0 Local Written Blocks 0 Femp Read Blocks 0			_totali,c.flag_totali,c.flag_hide_titolo,c.fla		
Local Dirtied Blocks				g_totall_sul_titolo,c.riag_conto		
Local Written Blocks			Index Cond	text) AND ((c.codice_template_fiscale)::te		
Temp Read Blocks				xt = (rc.codice_template_fiscale)::text) A ND (c.id_pdc = rc.id_pdc))		
Temp Written Blocks	0		Rows Removed by Index R	0		

The right node is the culprit

INDEX SCAN	546.93ms 5 %
on solving.conti (c)	8
Index Scan Node finds perform 2 read operation: actual value from the tabl	relevant records based on an Index . Index Scans s: one to read the index and another to read the le.
Node Type	Index Scan
Parent Relationship	Inner
Parallel Aware	false
Scan Direction	Forward
Index Name	cnti_pk
Relation Name	conti
Schema	solving
Alias	c
Startup Cost	0.29
Total Cost	0.35
Plan Rows	1
Plan Width	29
Actual Startup Time	0.019
Actual Total Time	0.037
Actual Rows	83
Actual Loops	14782
Output	c.id_conto,c.codice,c.descrizione,c.id_con to_parent,c.stato,c.tipo_mastro,c.ca,c.cli_ for,c.id_cee,c.id_riclass,c.tree_path,c.codi ce_template_fiscale,c.id_pdc,c.codice_azi enda,c.codice_tipo_sottoconto_anag,c.fla g_protetto,c.codice_part,c.flag_mostra_s empre_sottoconti,c.tipo_cespite,c.codice _stampa,c.flag_hide_codice,c.descrizione _totali,c.flag_totali,c.flag_hide_titolo,c.fla g_totali_sul_titolo,c.flag_conto
Index Cond	(((c.codice_part)::text = (rc.codice_part):: text) AND ((c.codice_template_fiscale)::te xt = (rc.codice_template_fiscale)::text) A ND (c.id_pdc = rc.id_pdc))

From the SQL query we expect the JOIN to the conti table to be performed after the join to the sottoconti table

BUT the optimizer decides to join conti to the result of the preceding join operations before joining to sottoconti

The node uses cnti_pk (PK index) but since it does not have the sottoconti row data it can only use three of the four primary key columns

The optimizer "thinks" this will return 1 row (Plan rows) but this is NOT true (83 rows returned)

This is a BAD result I can only explain by the planner assuming statistical independence on the index columns. In this case these 3 columns are perfectly correlated!

This causes the join to produce a cartesian product between the 14,782 rows of the left node and the 83 rows of this node (14,782 * 83 = 1,226,906)

How can we rewrite the SQL?

- Observe: conti and sottoconti are not relevant for filter and aggregation
- They are only used for projection
- We can commute the aggregation with the conti and sottoconti joins
 - So we will join to a smaller cardinality set
- This can be done introducing a subquery in the FROM clause (nesting)

The rewritten SQL

```
SELECT A0.*, Al.descrizione, Al.codice, Al.tree path,
(SELECT A2.tipo mastro FROM CONTI A2 WHERE A2.CODICE PART = A1.CODICE PART
AND A2.CODICE TEMPLATE FISCALE = A1.CODICE TEMPLATE FISCALE
AND A2.ID PDC = A1.ID PDC
AND A2.ID CONTO = A1.ID CONTO) AS tipo mastro
FROM (
 SELECT d.codice azienda
    , CASE
        WHEN d.id divisione IS NULL
       THEN '-'
       ELSE d.id divisione::text
      END || '#' || d.id esercizio comp || '#' || rc.id sottoconto AS record key
   , d.id divisione
    , d.id esercizio comp
    , sum (CASE
            WHEN rc.segno = 1
             AND rc.extra contabile = '0'
            AND coalesce(tda.sottotipo, '-') <> 'CH'
     THEN rc.importo
           ELSE 0
          END) AS totale dare
     , sum (CASE
            WHEN rc.seqno = (-1)
             AND rc.extra contabile = '0'
            AND coalesce(tda.sottotipo, '-') <> 'CH'
         THEN rc.importo
            ELSE 0
          END) AS totale avere
     , sum (CASE
            WHEN rc.seqno = 1
             AND rc.extra contabile = '0'
            AND coalesce(tda.sottotipo, '-') = 'CH'
         THEN rc.importo
            ELSE 0
          END) AS totale dare ch
    , sum (CASE
            WHEN rc.seqno = (-1)
            AND rc.extra contabile = '0'
            AND coalesce(tda.sottotipo, '-') = 'CH'
         THEN rc.importo
            ELSE 0
          END) AS totale avere ch
    , sum(CASE
            WHEN rc.seqno = 1
             AND rc.extra contabile = '1'
            AND coalesce(tda.sottotipo, '-') <> 'CH'
         THEN rc.importo
           ELSE 0
          END) AS totale_dare_ec
      CUP (CAPE
```

CREATE or replace VIEW saldi sottoconti AS

```
, sum(CASE
             WHEN rc.seqno = (-1)
               AND rc.extra contabile = '1'
               AND coalesce(tda.sottotipo, '-') <> 'CH'
             THEN rc.importo
             ELSE 0
           END) AS totale avere ec
     , rc.codice part
     , rc.codice template fiscale
     , rc.id pdc
     , rc.id sottoconto
FROM documenti d
join tipi doc az tda
  on tda.codice azienda = d.codice azienda
 and tda.codice tipo doc = d.codice tipo doc
 and tda.codice tipo doc az = d.codice tipo doc az
JOIN reg contabili rc
  ON d.codice azienda = rc.codice azienda
 AND d.id documento = rc.id documento
WHERE d.stato = 'C'
GROUP BY d.codice azienda
       , d.id esercizio comp
       , d.id divisione
       , rc.codice part
       , rc.codice template fiscale
       , rc.id pdc
       , rc.id sottoconto
) AS A0
JOIN SOTTOCONTI A1
ON A1.CODICE PART = A0.CODICE PART
AND A1.CODICE TEMPLATE FISCALE = A0.CODICE TEMPLATE FISCALE
AND A1.ID PDC = A0.ID PDC
AND A1.ID SOTTOCONTO = A0.ID SOTTOCONTO
```

The optimal plan



This plan is more than 76x faster than the original

The pattern I used

- Move joins used only for projection outside the aggregation
 - need to check the rewrite preserves semantics!
- Eventually force join order using subqueries in the SELECT clause
 - Only if 1 column selected
 - Otherwise introduce additional subquery level

What about WITH?

- While WITH could be used to bring the original query to 1.6s
- It can't be used in this VIEW since predicate pushdown does not happen for WITH (at the moment)
- The codice_azienda restriction cant' be hard coded!
- Without the restriction the query takes >16s which is worse than the original

Another example

REATE OR REPLACE VIEW intrastat 1c AS	JOIN righe_documento rd LEFT JOIN bolle_righe br2
ELECT m.codice azienda,	ON rd.codice_azienda = d.codice_azienda ON br2.codice_azienda = br1.codice_azienda
m.num rif,	AND rd.id documento = d.id documento AND br2.numero bolla = br1.numero bolla
d.codice part,	AND rd.tipo_omaggio IS NULL AND br2.numero riga padre = br1.numero riga
d.id anagrafica,	JOIN articoli varianti av AND br2.numero riga padre IS NOT NULL
date part('year', d.data comp) - 2000 AS anno cor	ON av.codice azienda = rd.codice azienda WHERE brl.numero riga padre IS NULL
CASE	AND av.id articolo = rd.id articolo GROUP BY brl.codice azienda, brl.id fattura, brl.numero riga
WHEN m.periodicita = 'M' THEN date part('month'	AND av.codice variante = rd.codice variante) br
WHEN m.periodicita = 'T' THEN date part('quarte	JOIN codici intrastat i ON br.codice azienda = rd.codice azienda
END AS periodo comp,	ON av.intra id = i.intra id AND br id fattura = rd id documento
COALESCE(i.codice, 'MANCA') AS codice intrastat,	JOIN articoli a AND br numero riga fattura = rd id riga documento
round(sum((rd.importo + rd.costo trasporto rip)	ON a.codice azienda = av.codice azienda WHERE m tipo riepilogo = 'C'
round(sum(br.mn1 + br.mn2), 0) AS massa netta,	AND a.id articolo = av.id articolo GPOUD BV m codice azienda m num rif. d codice part. d id apagra
round(sum(br.qta_um_sec), 0) AS qta_um_sec	JOIN tipi articolo ta
FROM mod_intra m	ON ta.codice azienda = a.codice azienda
JOIN documenti d ON m.codice_azienda = d.codice_az	AND ta.codice tipo articolo = a.codice tipo
CASE	AND ta.tipo = 'M'
WHEN m.periodicita = 'M' THEN (('01-' m.per	LEFT JOIN (
WHEN m.periodicita = 'T' THEN (('01-' ((m.p	SELECT brl.codice azienda,
END	brl.id fattura,
AND d.data_doc <	brl.numero riga fattura.
CASE	COALESCE (sum (
WHEN m.periodicita = 'M' THEN (('01-' (mod	CASE
CASE	WHEN brl.peso > 0 THEN brl.peso
WHEN m.periodo = 12 THEN m.anno + 1	WHEN av 1.peso unitario IS NOT NULL AND av 1.um peso =
ELSE m.anno	WHEN a l.codice um = 'KG' THEN brl.gta
END	WHEN brl.codice um acgven = 'KG' THEN brl.gta acguisto
WHEN m.periodicita = 'T' THEN (('01-' (mod	WHEN a l.codice um acquisto = 'KG' THEN brl.gta * av l.
CASE	END) 0) AS mp1
WHEN m.periodo = 4 THEN m.anno + 1	COALESCE (sum (
ELSE m.anno	CASE
END	WHEN $br_{2,peso} > 0$ THEN $br_{2,peso}$
END	WHEN br2.codice um acquer = 'KG' THEN br2.gta acquisto
JOIN tipi_doc_az taz	END), 0) AS mp2.
ON taz.codice_azienda = d.codice_azienda	sum (
AND taz.codice_tipo_doc = d.codice_tipo_doc	
AND taz.codice_tipo_doc_az = d.codice_tipo_doc_a	WHEN a 1 codice um = i 1 codice um sec THEN brl gta
AND taz.codice_tipo_doc = 'VEN'	WHEN brl codice um acquer = i l codice um sec THEN brl
JOIN numeratori nu	WHEN a l codice um acquieta = i l codice um sec THEN by
ON nu.codice_azienda = d.codice_azienda	FND) AS gta um seg
AND nu.anno = d.anno	FPOM balle righe bri
AND nu.id_numeratore = d.id_numeratore	TOIN articoli varianti av 1
AND nu.tipo_protocollo <> 2	ON av 1 godige avierda - brl godige avierda
JOIN anagrafiche_vr vr	AND av 1 id articolo = brl id articolo
ON vr.codice_part = d.codice_part	AND av 1. redice variante - bri codice variante
AND vr.id_anagrafica = d.id_anagrafica	TOTM articoli a 1
AND vr.prog_vr = d.prog_vr_anag	ON all addies priords - av 1 addies priords
JOIN nazioni n	AND a lid atticule - av lid atticule
ON n.codice_nazione = vr.codice_nazione	NUD a_1.10_atticolo = av_1.10_atticolo
AND n.cee = '1'	
AND n.codice_nazione <> 'ITA'	ON av_1.intra_ld = 1_1.intra_ld

The query

SELECT *
FROM intrastat_1c
WHERE codice_azienda = 'S0032IC'
AND num_rif = 451

This is the PRIMARY KEY of the mod_intra table

The plan



This is the slow plan. Takes 1 min and 16 seconds.

Let's zoom in



Continue zooming...

INDEX SCAN on solving.bolle righe (br2) 59.97s | **86** %

using bolr_pk

Index Scan Node finds relevant records based on an Index. Index Scans perform 2 read operations: one to read the index and another to read the actual value from the table.

Node Type	Index Scan
Parent Relationship	Inner
Parallel Aware	false
Scan Direction	Forward
Index Name	bolr_pk
Relation Name	bolle_righe
Schema	solving
Alias	br2
Startup Cost	0.42
Total Cost	331.04
Plan Rows	1
Plan Width	34
Actual Startup Time	2.369
Actual Total Time	2.405
Actual Rows	0
Actual Loops	24936
Index Cond	(((br2.codice_azienda)::text = (br1.codice _azienda)::text) AND ((br2.codice_aziend a)::text = 'S0032IC'::text) AND ((br2.num ero_bolla)::text = (br1.numero_bolla)::te xt))
Rows Removed by Index R echeck	0
Filter	((br2.numero_riga_padre IS NOT NULL) AND (br2.numero_riga_padre = br1.num ero_riga))
Rows Removed by Filter	10
Shared Hit Blocks	4242705

Here the developer had the join wrong and forgot to add a condition. This causes the index to be used inefficiently (see large value of Shared Hit Blocks).

So we fix the query adding the missing clause.

This is MUCH better



This fixed version talkes 7.83 seconds. Are we done?

We can observe that the query has two main branches and the right one depicted here actually does not depend on the restriction criteria we have in the query. PostgreSQL does its best to optimize this subquery but here we have a case of linear complexity where the query time is going to grow steadily as more data is added to the bolle_righe table. This is going to be a problem in the future assuming the current performance result is acceptable.

We should rewrite the query so that the right branch can be bound in size depending on the results from the left branch which actually depends on the query restriction.

The strategy

- Rewrite the query to force a particular join order
- This can be done using the nesting trick, using subqueries in the FROM clause to our advantage
- The rewritten query is more complex sintactically and less readable, but is much faster, as we will see

The new VIEW definition

CREATE OR REPLACE VIEW intrastat 1c AS SELECT codice azienda, num rif, codice part, id anagrafica, anno comp, periodo comp, codice intrastat, codice um sec, ROUND(SUM(ammontare), 0) AS ammontare, ROUND(SUM(peso), 0) AS massa netta, ROUND(SUM(qta um sec), 0) AS qta um sec FROM (SELECT tl.codice azienda, tl.num rif, tl.id documento, tl.id riga documento, tl.codice part, tl.id anagrafica, tl.anno comp, tl.periodo comp, i.codice AS codice intrastat, i.codice um sec, MIN(t1.ammontare) AS ammontare, sum (COALESCE (CASE WHEN brl.peso > 0 THEN brl.peso WHEN av.peso unitario IS NOT NULL AND av.um peso = 'KG' THEN av.peso unitario * WHEN a.codice um = 'KG' THEN brl.gta WHEN brl.codice um acqven = 'KG' THEN brl.gta a WHEN a.codice um acquisto = 'KG' THEN brl.qta * END, CASE WHEN br2.peso > 0 THEN br2.peso WHEN br2.codice um acqven = 'KG' THEN br2.gta a END, 0)) AS peso, sum (COALESCE (CASE WHEN a.codice um = i.codice um sec THEN brl.qta WHEN brl.codice_um_acqven = i.codice_um_sec THE WHEN a.codice um acquisto = i.codice um sec THE END, 0)) AS qta um sec FROM (SELECT m.codice azienda, m.num rif, d.id documento, rd.id riga documento, d.codice part, d.id anagrafica, date part('year', d.data comp) - 2000::double CASE WHEN m.periodicita = 'M' THEN date part('mor WHEN m.periodicita = 'T' THEN date part('qua

END AS periodo comp. rd.importo + rd.costo trasporto rip * taz.segno 1 FROM mod intra m JOIN documenti d ON m.codice azienda = d.codice azienda AND d.data doc >= CASE WHEN m.periodicita = 'M' THEN (('01-' || m.periodo) WHEN m.periodicita = 'T' THEN (('01-' || ((m.periodo END::date AND d.data doc < CASE WHEN m.periodicita = 'M' THEN (('01-' || (mod(m.period CASE WHEN m.periodo = 12 THEN m.anno + 1 ELSE m.anno END WHEN m.periodicita = 'T' THEN (('01-' || (mod(m.perioc CASE WHEN m.periodo = 4 THEN m.anno + 1 ELSE m.anno END END::date JOIN tipi doc az taz ON taz.codice azienda = d.codice azienda AND taz.codice tipo doc = d.codice tipo doc AND taz.codice tipo doc az = d.codice tipo doc az AND taz.codice tipo doc = 'VEN' JOIN numeratori nu ON nu.codice azienda = d.codice azienda AND nu.anno = d.anno AND nu.id numeratore = d.id numeratore AND nu.tipo protocollo <> 2 JOIN anagrafiche vr vr ON vr.codice part = d.codice part AND vr.id anagrafica = d.id anagrafica AND vr.prog vr = d.prog vr anag JOIN nazioni n ON n.codice nazione = vr.codice nazione AND n.cee = '1' AND n.codice nazione <> 'ITA' JOIN righe documento rd ON rd.codice azienda = d.codice azienda AND rd.id documento = d.id documento AND rd.tipo omaggio IS NULL WHERE m.tip_riepilogo = 'C') AS t1 JOIN bolle righe brl ON brl.codice azienda = tl.codice azienda AND brl.id fattura = tl.id documento AND brl.numero_riga_fattura = tl.id_riga_documento JOIN articoli_varianti av ON av.codice azienda = brl.codice azienda AND av.id articolo = brl.id articolo AND av.codice variante = brl.codice variante

JOIN articoli a ON a.codice azienda = av.codice azienda AND a.id articolo = av.id articolo JOIN codici intrastat i ON av.intra id = i.intra id LEFT JOIN bolle righe br2 ON br2.codice azienda = br1.codice azienda AND br2.codice magazzino = br1.codice magazzino AND br2.numero bolla = br1.numero bolla AND br2.numero riga padre = br1.numero riga WHERE brl.numero riga padre IS NULL AND brl.id fattura IS NOT NULL GROUP BY tl.codice azienda, tl.num rif, tl.id documento, tl.id riga documento, tl.codice part, tl.id anagrafica, tl.anno comp, tl.periodo comp, i.codice, i.codice um sec) AS RESULT GROUP BY codice azienda, num rif, codice part, id anagrafica, anno comp, periodo comp, codice intrastat, codice um sec

We have two GROUP BY clauses in order to get amounts correct. We have one amount but multiple weights coming from bolle_righe. A single GROUP BY would sum amounts multiple times.

The OPTIMAL plan



This plan takes 19ms for execution + 24ms for planning. This is 178x ;-) faster than the previous solution and, even more importantly, <u>it maintains constant performance</u> as the data size grows.

So, to RDBMS deterrents we can say that a 12-way join can take a few milliseconds...;-)

Summary and thoughts

- Perform these optimizations only when required
- Optimal SQL sometimes needs to be written in a less declarative way
- For views, beware of multiple use cases. You may optimize one but worsen others
- Think about expected data flows and check if the plan matches. Many times the planner does better than you, but sometimes it gets things wrong
- Don't be afraid of using nesting when it can be useful. PostgreSQL is good at pushing down predicates (WITH is currently an exception)
- Think about branch complexity: avoid linear complexity (branches that don't take advantage of query restrictions)
- Could PostgreSQL do these rewrites by himself?